

Large Scale Data Management

Graph Stores, Neo4J

Silviu Maniu, LIG, Univ. Grenoble Alpes

Graphs

Graphs correspond to a natural organization of knowledge

They generalize

- Relations
- Trees (documents)
- · Key-value pairs

Graph stores simplify / facilitate data representation

They do not simplify query evaluation (and may make it more complex)

Graph Data

Graphs

Data model: *G(V,E)* edges *E* between vertices *V*

Edges can be: directed vs. undirected, weighted vs. unweighted

Nodes/vertices can be:

- Unlabelled
- Having a single label or type
- Set of attribute-value pairs (property graphs)

Graph can have semantics (RDF, knowledge graphs)

Example Graphs

Social Networks

Edges: persons, nodes: relationships

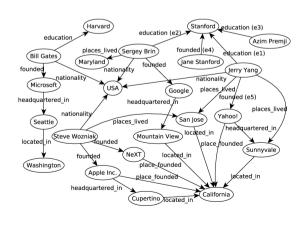


by Michael Coghlan, CC BY-SA 2.0 via Flickr

Example Graphs

Knowledge Graphs / Semantic Data

Edges: subjects/objects, nodes: predicates



Example Graphs

Road Networks

Edges: roads/connections, nodes: intersections/cities/etc.



Graph Data ModelsRDF

RDF: Data model for (originally) data interchange on the Web

Structured in triples: (subject, predicate, object)

Graph: triples form a directed, labeled/typed, graph

Nodes: are generally URIs allowing to address everything

Schemas: RDF schema (RDFS), OWL



Graph Data Models

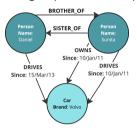
Property Graphs

Nodes: entities / objects / concepts

Edges: directed relationships between nodes

Properties: key-values pairs to nodes/edges (name: "Alice")

Labels: edge/node types



Graphs and Algorithms

Distributing Graph Computation

Challenges in efficiently processing graph algorithms:

- locality issues: since different vertices are needed
- little work per vertex: some vertices may not be used at all
- parallelism is not the same everywhere

MapReduce, Spark are data parallel: very efficient for tabular data, aggregating data – not the case for graphs

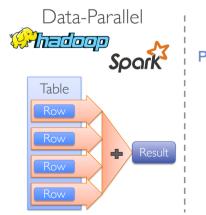
Graphs and Algorithms

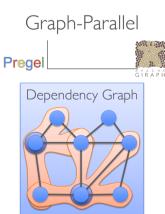
Graphs and algorithms on graphs are important:

- classic algorithms: shortest paths, community detection, counting triangles, PageRank
- data mining/machine learning on graphs: collaborative filtering, belief propagation, mean-field algorithms, graph neural networks
- · querying: matching patterns / paths on graphs

Graph Parallel vs. Data Parallel

Graph Parallel: computation in each vertex





Graph Processing with Pregel

Implements a variant of the Bulk Synchronous Processing model 6 6 2 Supersto

Computation Model

- 1. input: each vertex in the graph receives initial data
- 2. supersteps with global synchronization points
 - each vertex receives the same user-defined function to execute in a superstep
 - between supersteps, messages are sent between vertices (either via links, or globally)
 - in a superstep, each vertex can modify its state, and set-up a message to be sent to other vertices
 - vertices can vote to halt the computation; if they receive new messages they must re-activate
- 3. computation stops when all vertices vote to halt and are not reactivated

Neo4J Model

Property Graphs

Neo4J uses a property graph data model (directed, labeled, key-value properties)



Graph Stores: Neo4J

Neo4J Model

Data Manipulation Language

Data manipulation language (CRUD): **Cypher**, used to describe data and patterns to be matched Node descriptions in Cypher:

Neo4J Model

Data Manipulation Language

```
Relationship descriptions in Cypher:

-- (undirected) vs. --> or <-- (directed)

Sample relationship descriptions:

-->

-[role]-> - relationship ID

-[:ACTED_IN]-> - relationship type

-[role:ACTED_IN]->

-[role:ACTED_IN {roles: ["Neo"]}]->

-relationship with attributes
```

Neo4J Architecture

Transaction Control, Clustering

Transaction control is ACID when one server

• Write-ahead logs (WAL) for recovery - records the transaction log

Distribution/clustering, servers dealing with one or more databases:

- Primary mode: for safety and transaction control
- Secondary mode: for scalability and read performance

When clustering

 Causal consistency: each client writing to a primary server is guaranteed to read the same data from a secondary server

Neo4J Architecture

Database Organisation

On-disk database in files:

- nodestore node related data
- relationship relationship related data
- property properties (key/value)
- label (for indexes) index related label data

Schema-less: fixed record size for quick random access (via offsets), stored on disk but cached in memory

Optimized for graph traversal = linked list, every node and relationship has pointers to the first in the chain

Data Manipulation with Cypher

Patterns combine node and relationship descriptors:

```
(keanu:Person:Actor {name: "Keanu Reeves"} )
-[role:ACTED_IN {roles: ["Neo"] } ]->
(matrix:Movie {title: "The Matrix"} )

Data creation:

CREATE (a:Person { name: "Tom Hanks", born:1956 })
-[r:ACTED_IN { roles: ["Forrest"]}]-> (m:Movie { title: "Forrest Gump", released:1994 })

CREATE (d:Person { name: "Robert Zemeckis", born:1951 }) -[:DIRECTED]->(m)
```

Querying with Cypher

MATCH pattern RETURN matched variables

Successive match-create-return steps can be used to update the data:

```
MATCH (p:Person { name:"Tom Hanks" })
CREATE (m:Movie { title:"Cloud
Atlas",released:2012 })
CREATE (p)-[r:ACTED_IN { roles: ['Zachry']}]->(m)
RETURN p,r,m
```

Querying with Cypher

Returning Results

```
MATCH (a { name: "A" })-[r]->(b)
RETURN * //all results, all properties of a,r,b

MATCH (n)
RETURN n.age // returns null if no age

MATCH (a { name: "A" })
RETURN a.age > 30, "I'm a literal", (a)->() //edge creation
```

Querying with Cypher

```
Inserting data only if it didn't exist:
```

```
MERGE (m:Movie { title:"Cloud Atlas" })
   // create or check the existence of movie node m
ON CREATE SET m.released = 2012
   // if we had to create it, set the release year

RETURN m
Insert relationship only if it did not exist:

MATCH (m:Movie { title:"Cloud Atlas" })
MATCH (p:Person { name:"Tom Hanks" })
MERGE (p)-[r:ACTED IN]->(m)
ON CREATE SET r.roles =['Zachry']
RETURN p,r,m
```

Other Cypher Functionalities

· Boolean conditions:

```
MATCH (n)

WHERE n.name = 'Peter' XOR (n.age < 30 AND n.name = "Tobias") OR
NOT (n.name = "Tobias" OR n.name="Peter")</pre>
```

- RETURN n
- Optional matching: somewhat like outer joins, replace missing with ${\tt NULL}$
- Returned data can be: ordered, truncated, aggregated
- Unwind: unfolds a collection into a set UNWIND[1,2,3] AS x RETURN x // three results
- Indexes: CREATE INDEX ON : Person (name)
- EXPLAIN to get the query plan, PROFILE to measure the effort

Neo4J Architecture

Query Perfomance via Indexes

Search performance (automatically used if present):

- Range: default, B+ trees
- Text and Point: for string and coordinate operations
- Lookup: only for node label and relationship type (activated by default)

Semantic indexes (not automatically used):

- Full-text: search for content, via Apache Lucene
- Vector: similarity search via embeddings

To Read Further

Articles

- Malewicz et al., Pregel: a system for large-scale graph processing, SIGMOD 2010
- Rodriguez, Neubauer, Constructions from Dots and Lines, Bul. Am. Soc. Info. Sci. Tech. 36(6), 2010
- Francis et al., Cypher: An Evolving Query Language for Property Graphs, SIGMOD 2018

Documentation

Neo4J: https://neo4j.com/docs/ cypher/

Summary

Query Perfomance via Indexes

- Very convenient data model, natural representation
- · Typically no strict schema
- No standard query language (effort to extend Cypher to standard, GQL)
- Knowledge graphs are a particular case (RDF and SPARQL are standards)
- Other powerful tools around: distributed graph stores (Pregel, Spark GraphX)
 - · Extra dimension: graph partitioning
 - · Less effort on query language; in progress